

General radix, recursive Cooley-Tukey FFT

assume $n = km$

$$\text{DFT}_{km} = (\underbrace{\text{DFT}_k}_{\text{radix}} \otimes \underbrace{I_m}_{\text{diagonal matrix}}) \underbrace{T_m}_{\text{diagonal matrix}} (\underbrace{I_k}_{\text{diagonal matrix}} \otimes \text{DFT}_m) L_k^n$$

3 key structures: $I_k \otimes A_m$, $A_k \otimes I_m$, L_k^n

1.) $Y = (I_k \otimes A_m) X$

$$\begin{pmatrix} Y \\ Y \\ \vdots \end{pmatrix} = \begin{pmatrix} A & & \\ & A & \\ & & \ddots \\ & & & A \end{pmatrix} \begin{pmatrix} X \\ X \\ \vdots \end{pmatrix}$$

k A's at stride 1

for $i = 0 : k-1$
 $Y[i:m:i+m-1] = A \cdot X[i:m:i+m-1]$

2.) $Y = (A_k \otimes I_m) X$

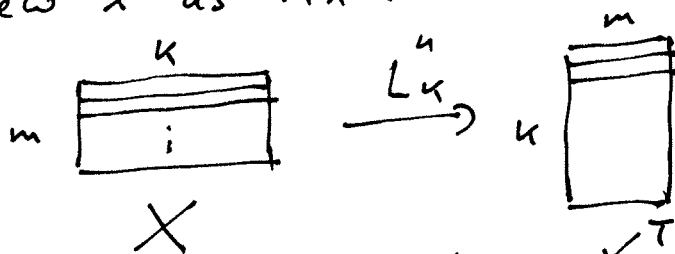
$$\begin{pmatrix} Y \\ Y \\ \vdots \end{pmatrix} = \begin{pmatrix} A & & & \\ & A & & \\ & & \ddots & \\ & & & A \end{pmatrix} \begin{pmatrix} X \\ X \\ \vdots \end{pmatrix}$$

m A's at stride m

for $i = 0 : m-1$
 $Y[i:m:i+(k-1)m] = A \cdot X[i:m:i+(k-1)m]$

3.) $Y = L_k^n X$: different ways of viewing it

a.) view X as $m \times k$ matrix:



transposition!

b.) L_k^n "reads at stride 1 and writes at stride m "

stride 1 $\rightarrow m$
 is the same as
 stride $k \rightarrow 1$

c.) L_k^n performs permutation $im+j \rightarrow jk+i$ $0 \leq i < k$
 $0 \leq j < m$

FFT again:

$$\text{DFT}_{km} = (\underbrace{\text{DFT}_k \otimes I_m}_{\text{stride } m \rightarrow m}) \underbrace{T_m}_{\text{stride } 1 \rightarrow 1} (\underbrace{I_k \otimes \text{DFT}_m}_{\text{stride } 1 \rightarrow m}) \underbrace{L_k^n}_{\text{stride } 1 \rightarrow m}$$

stride 1 $\rightarrow m$
 is the same as
 stride $k \rightarrow 1$

this is the "decimation-in-time" version

Decimation in frequency: transpose:

Use: - DFT is symmetric

$$- (L_n^u)^T = L_m$$

$$- (A \otimes B)^T = A^T \otimes B^T$$

Gives:

$$\text{DFT}_{kn} = L_m^u (I_k \otimes \text{DFT}_m) T_m^u (\text{DFT}_k \otimes I_m)$$

Cost analysis: (was in exam for radix 2), assume $n=2$

Measure: (complex adds, complex mults)

Cost: independent of radix $(n \log_2(n), \frac{1}{2} n \log_2(n))$

$$\begin{array}{lcl} \text{complex add} & = & 2 \text{ real adds} \\ \text{" mult} & \leq & 4 \text{ real mults} \\ & & 2 \text{ real adds} \end{array}$$

$$\Rightarrow \text{real cost} \leq 2n \log_2(n) + 3n \log_2(n) = 5n \log_2(n)$$

Iterative radix-2 FFT

$$\text{DFT}_{2^t} = R_{2^t} \prod_{i=1}^t \underbrace{(I_{2^{t-i}} \otimes T_{2^{t-i}}^{2^i})}_{\text{diagonal matrix}} \underbrace{(I_{2^{t-i}} \otimes \text{DFT}_2 \otimes I_{2^{t-i}})}_{2^{t-i} \text{ DFT}_2\text{'s at varying strides}$$

$\underbrace{\hspace{10em}}_{\text{bit-reversal permutation}}$

Most people consider this "the FFT"