

A presentation by Samuel Häusler about the paper:

SPL

A Language and Compiler for DSP Algorithms

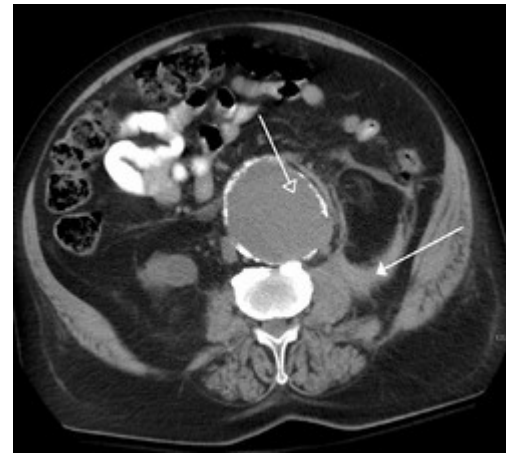
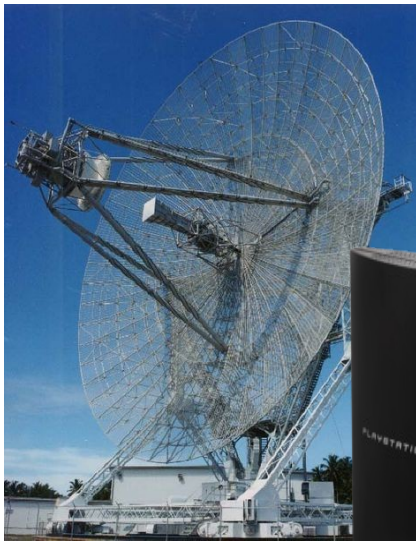
Jiangxing Jong, Jeremy Johnson, Robert Johnson and David Padua

Overview

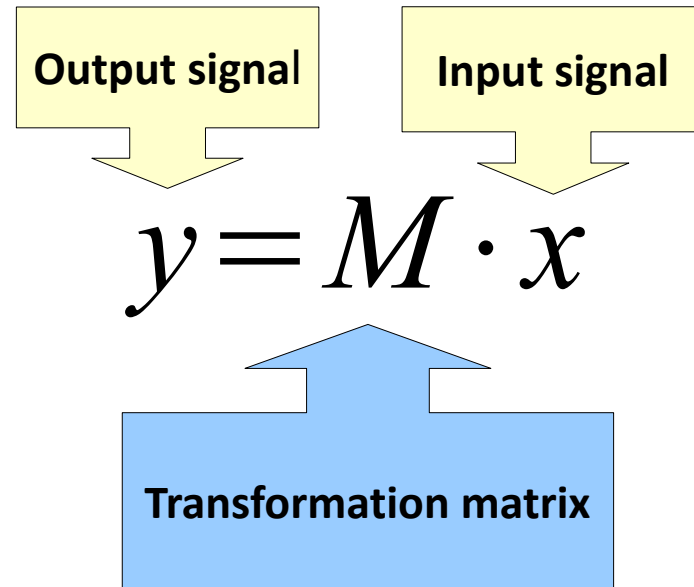
- **Motivation**
- **The language SPL**
- **The SPL compiler**
- **Experiments and results**
- **Conclusion**

Motivation

Digital Signal Processing is part of everyday life



Signal transformations:



Example: Discrete Fourier transformation (1D)

$$y = F_n \cdot x \quad \text{where } F_{pq} = \omega_n^{pq} \quad \text{with } \omega = e^{\frac{2\pi i}{n}}$$

$$F_4 = \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & -i & -1 & i \\ 1 & -1 & 1 & -1 \\ 1 & i & -1 & -i \end{pmatrix}$$

Fast signal transforms

$$\left| \begin{array}{c} x_1 \\ x_2 \\ x_3 \end{array} \right.$$

$$\overline{\begin{pmatrix} 1 & 0 & 2 \\ 0 & 3 & 0 \\ 1 & 0 & 2 \end{pmatrix}} \begin{pmatrix} x_1 + 2x_3 \\ 3x_2 \\ x_1 + 2x_3 \end{pmatrix}$$

$$\begin{pmatrix} 1 & 2 & 3 \\ 1 & 2 & 3 \\ 1 & 2 & 3 \end{pmatrix}$$

$$\left| \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} \right.$$

 y

$$\begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix} \begin{pmatrix} 1 & 2 & 3 \end{pmatrix}$$

$$\left| \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} \right.$$

 y

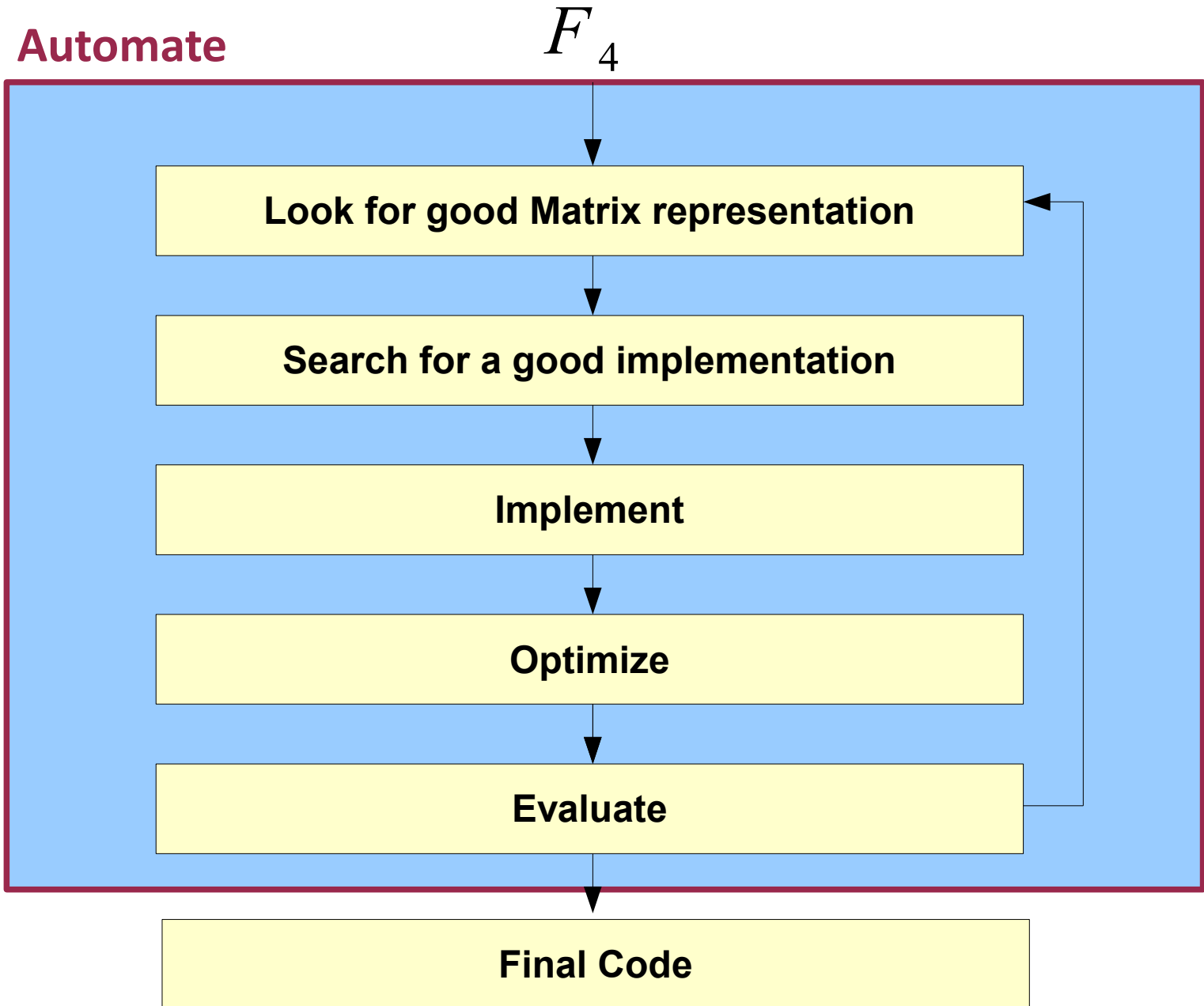
Fast Fourier Transform

$$F_4 = \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & -i & -1 & i \\ 1 & -1 & 1 & -1 \\ 1 & i & -1 & -i \end{pmatrix} = \begin{pmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & -1 & 0 \\ 0 & 1 & 0 & -1 \end{pmatrix} \cdot \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -i \end{pmatrix} \cdot \begin{pmatrix} 1 & 1 & 0 & 0 \\ 1 & -1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & -1 \end{pmatrix} \cdot \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

$$F_4 = (F_2 \otimes I_2) T_2^4 (I_2 \otimes F_2) L_2^4$$

$$F_2 = \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \quad I_2 = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$

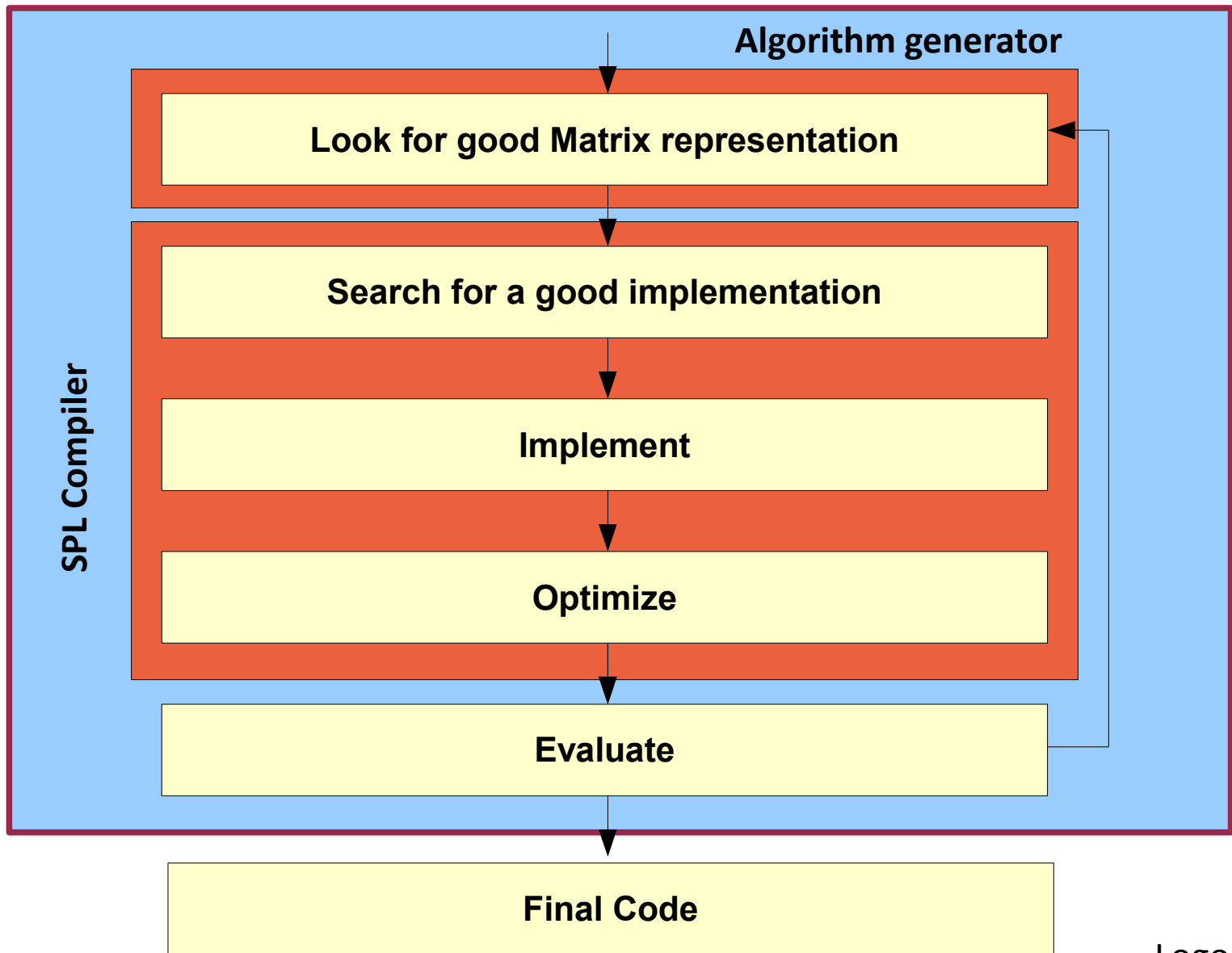
Implementation of signal transformations



Example:



F_4



Algorithm generator:

$$F_4 = (I_2 \otimes F_2) T_2^4 (I_2 \otimes F_2) L_2^4$$

```
(compose
  (tensor (F 2) (I 2))
  (T 4 2)
  (tensor (I 2) (F 2))
  (L 4 2))
```

Compiler

```
(template (compose A B))  
  
  [A.in_size==B.out_size]  
  
  (B_ ($in, $t0, 0, 0, 1, 1)  
A_ ($t0, $out, 0, 0, 1, 1))
```

Overview

- **Motivation**
- **The language SPL**
- **The SPL compiler**
- **Experiments and results**
- **Conclusion**

The language SPL

A SPL program is a transformation matrix

$y =$

```
#language c
(define A (I 1))
(define B (F 2))
#subname myfunction
(compose A B)
```

#Compiler directive
(matrix definitions)
 $\cdot x$
(matrix operation)

General matrix definitions

$$\text{(matrix ((a11 a12 a13) (a21 a22 a23) (a31 a32 a33))}$$

$$\begin{pmatrix} a11 & a12 & a13 \\ a21 & a22 & a23 \\ a31 & a32 & a33 \end{pmatrix}$$

$$\text{(diagonal (a11 a22 a33))}$$

$$\begin{pmatrix} a11 & 0 & 0 \\ 0 & a22 & 0 \\ 0 & 0 & a33 \end{pmatrix}$$

$$\text{(permutation (3 1 2))}$$

$$\begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{pmatrix}$$

Parametrized matrices

Identity matrix: $(I \ n)$

Fourier transformation: $(F \ n)$

Stride permutation matrix: $(L \ mn \ n)$

for index j : $F_r^n: j \Rightarrow j \cdot r \bmod n$

$$L_2^4 = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

Twiddle matrix: $(T \ mn \ n)$

$$T_m^{mn} = \begin{pmatrix} I & 0 & 0 \\ 0 & W_m & 0 \\ 0 & 0 & W_m^{(n-1)} \end{pmatrix}$$

Matrix operations

Matrix product: (compose A B)

Direct sum: (direct-sum A B)

Tensor product: (tensor A B)

$$\begin{pmatrix} 1 & 0 & 2 \\ 0 & 3 & 0 \\ 1 & 0 & 2 \end{pmatrix} \otimes A = \begin{pmatrix} A & 0 & 2A \\ 0 & 3A & 0 \\ A & 0 & 2A \end{pmatrix}$$

Name assignment:

```
(define name formula)
```

e.g. `(define A (I 1))`

Program example

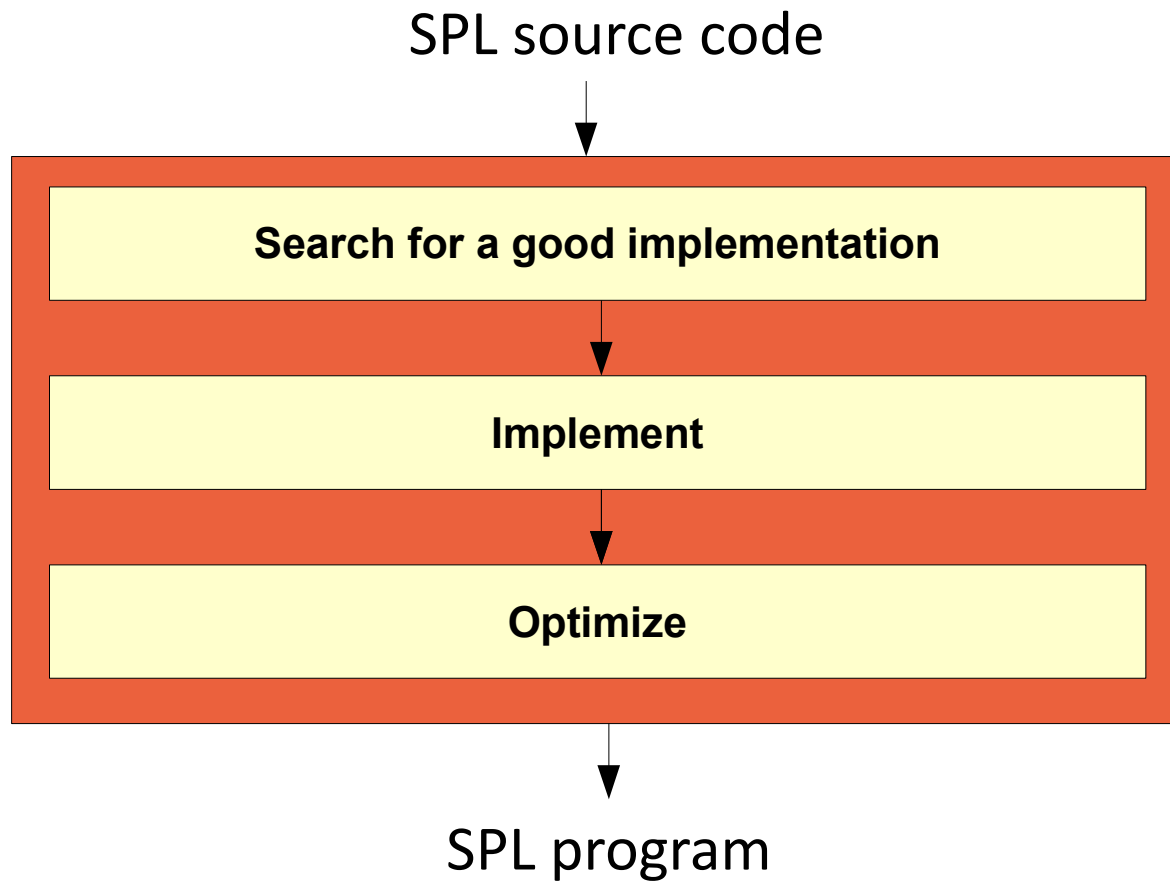
```
(define F4
  (compose
    (tensor (F 2) (I 2)) (T 4 2)
    (tensor (I 2) (F 2)) (L 4 2)))
#subname fft16
(compose
  (tensor F4 (I 4)) (T 16 4)
  (tensor (I 4) F4) (L 16 4))
```

$$F_{16} = (F_4 \otimes I_4) T_4^{16} (I_4 \otimes F_4) L_4^{16}, \text{ where } F_4 = (F_2 \otimes I_2) T_2^4 (I_2 \otimes F_2) L_2^4$$

Overview

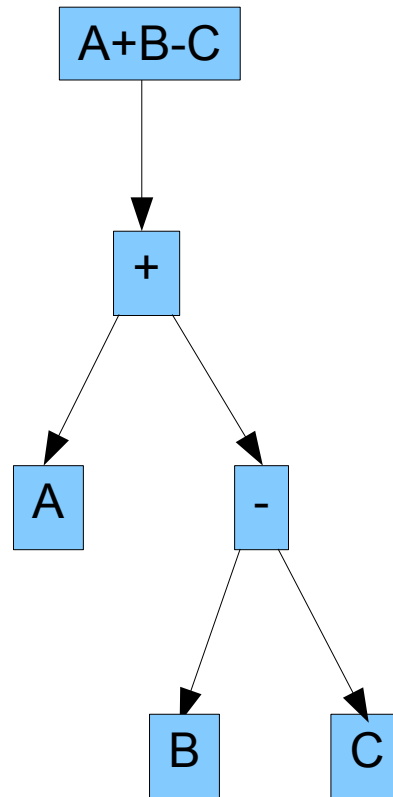
- **Motivation**
- **The language SPL**
- **The SPL compiler**
- **Experiments and results**
- **Conclusion**

The compiler

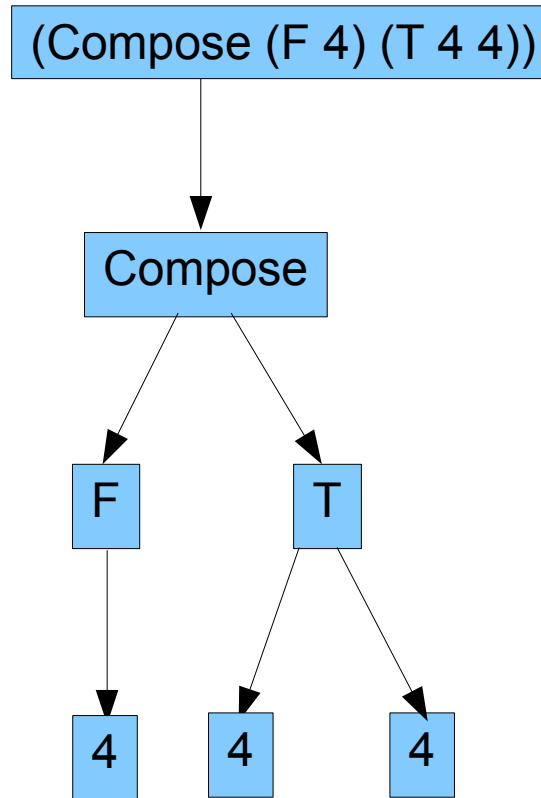


Parsing

Abstract syntax tree:



Example (compose (F 4) (T 4 4))



Templates

```
(template (F n_) [n_>0]
  (do $i=0, n_-1
    $out($i0)=0

    do $i1=0, n_-1
      $r0 = $i0 * $i1
      $f0 = W(n_ $r0)*$in($i1)
      $out($i0)=$out($i0)+$f0
    end
  end) )
```

Templates

```
(template pattern condition i-code)
```


Templates

Templates:

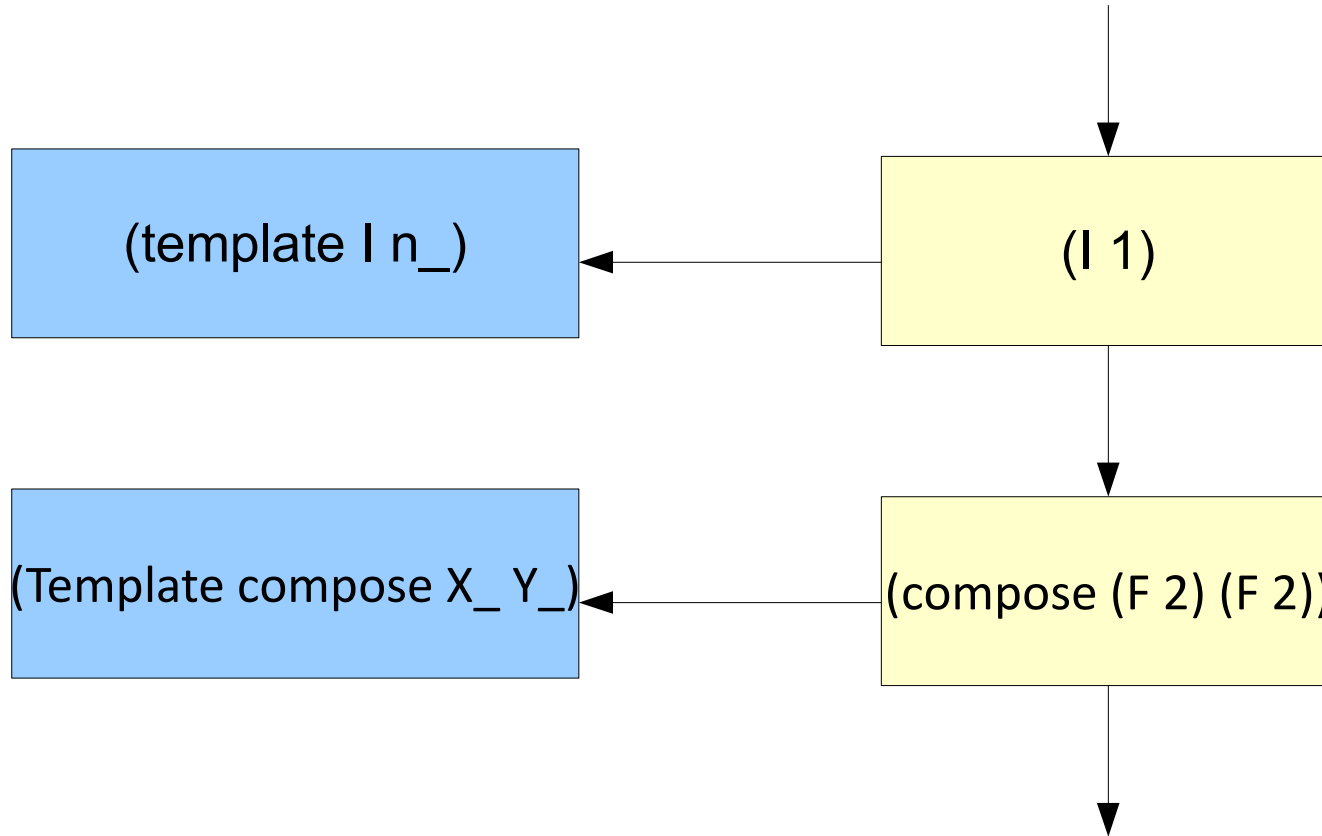
Compiler:

(template I n_)

(Template compose X_ Y_)

(I 1)

(compose (F 2) (F 2))



Templates

(L m_ n_) with condition [m_ == 2 * n]

(L 4 2)

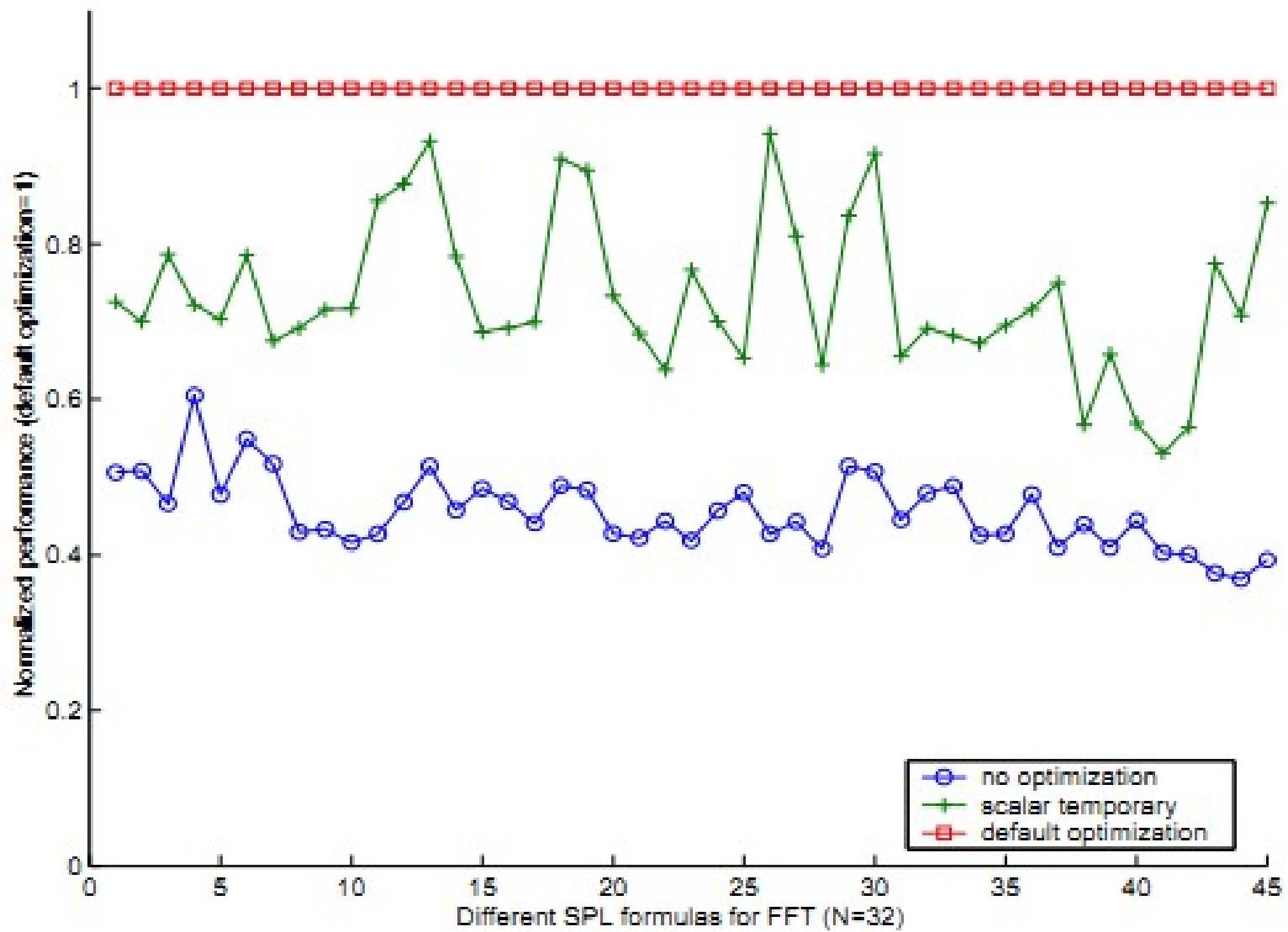
(L 4 1)

Templates

Example: A template for (F n)

```
(template (F n_) [n_>0]
  (do $i=0, n_-1
    $out($i0)=0

    do $i1=0, n_-1
      $r0 = $i0 * $i1
      $f0 = W(n_ $r0)*$in($i1)
      $out($i0)=$out($i0)+$f0
    end
  end) )
```



Overview

- **Motivation**
- **The language SPL**
- **The SPL compiler**
- **Experiments and results**
- **Conclusion**

Experiments and results

... on FFT transforms F_{2^i} for $i=1-20$

SPL Compiler

Search Strategy

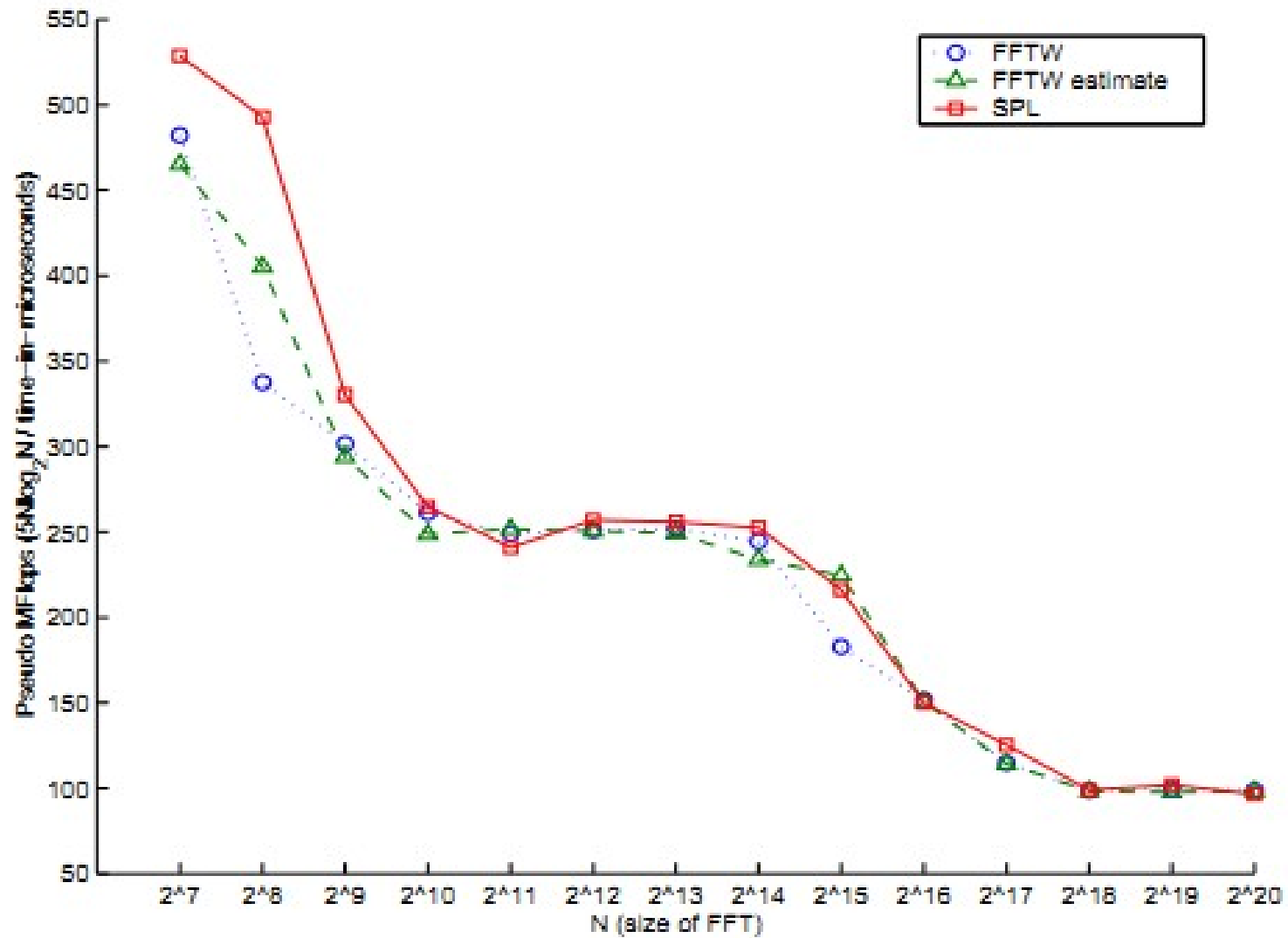
Compared results with „fftw“ library results

Search strategy

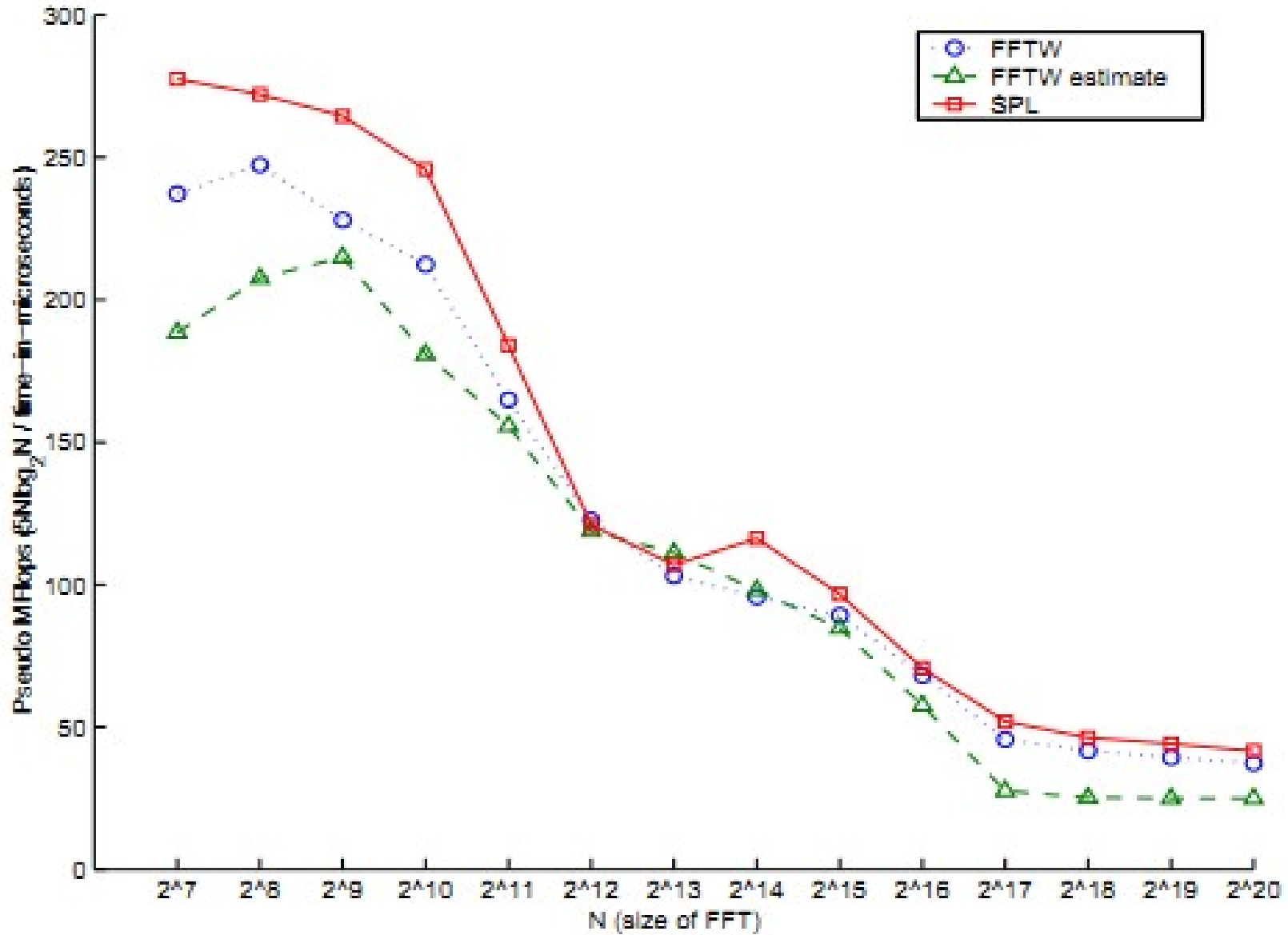
Dynamic programming over factorization by:

$$F_{rs} = (F_r \otimes I_s) T_s^{rs} (I_r \otimes F_s) L_r^{rs}$$

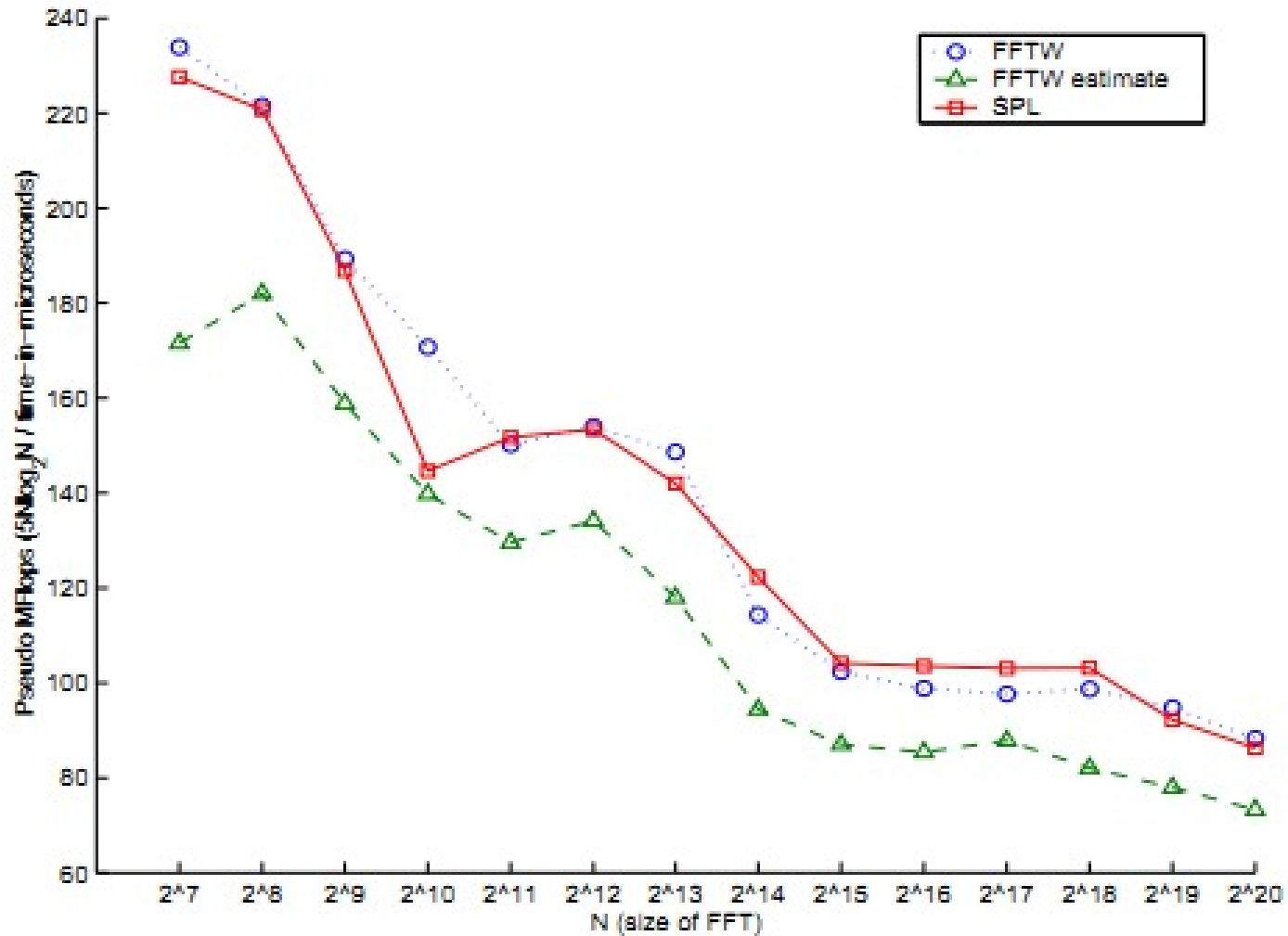
SPARC



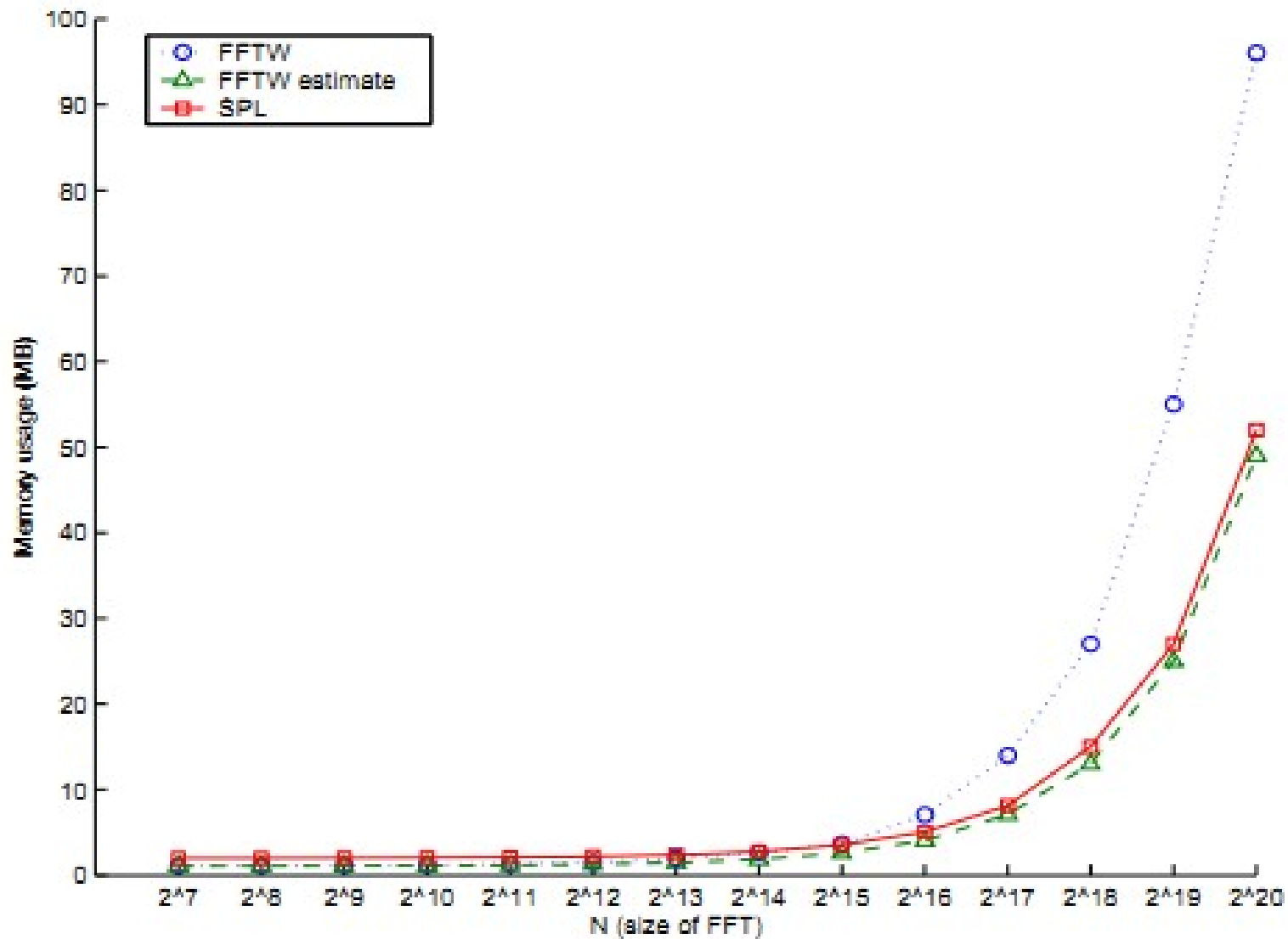
MIPS



Pentium II



Memory usage



Conclusions

- Competitive with fftw if connected with search library
- More general, not only fft
- Results show power of domain-specific language and compilers
- Many more domain-specific projects in future