

Fast Transform Algorithms via Projection on a Subalgebra

Yevgen Voronenko

Frédéric de Mesmay

Electrical and Computer Engineering
Carnegie Mellon University

Introduction

- Fast $n \log n$ algorithms are important for practical applications
- Deriving such algorithms is typically tedious and not well understood process
- We present a general method for mechanical derivation of new algorithms by *projecting* existing algorithms
- We successfully obtain fast algorithms for all 16 types of DCTs/DSTs

Related Work : Trigonometric Transform Algorithms

- **Hong and Vetterli, 1991**
 - “Basefield Transforms with Convolution Property”
 - Algebraic projection (by human)
- **Frigo and Johnson, 1999**
 - “Fast Fourier Transform Compiler”
 - Data-flow graph projection (by a computer)
- **Pueschel and Egner, 2002**
 - Symmetry Based Matrix Factorization
 - Automatic algorithm derivation from a matrix (by a computer)
- **Pueschel and Moura, 2001—2006**
 - Algebraic Theory of Signal Processing
 - Cooley-Tukey type algorithms for all 16 kinds of DCTs/DSTs

**ASP allows to understand what's going on
and thus mechanically derive algorithms**

Outline

1. Introduction
2. Embedding method
3. Algorithm projection method

Embedding Method: Problem statement

■ Given:

- Regular signal model (\mathcal{A}, Φ_A) + Fourier transform \mathcal{F}_A
- Regular signal model (\mathcal{B}, Φ_B) + Fourier transform \mathcal{F}_B
- $\mathcal{B} \leq \mathcal{A}$

■ Find:

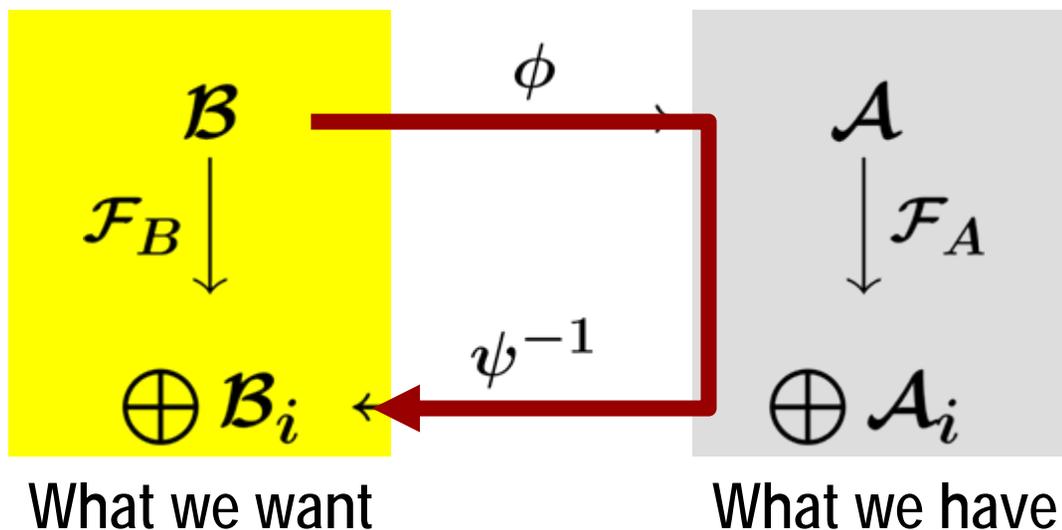
- Express \mathcal{F}_B using \mathcal{F}_A

Embedding Method

- Basic idea: embed \mathcal{B} in \mathcal{A}
- Let ϕ and Ψ be embeddings

$$\phi : \mathcal{B} \rightarrow \mathcal{A} \qquad \psi : \bigoplus \mathcal{B}_i \rightarrow \bigoplus \mathcal{A}_i$$

- Following diagram commutes



Matrix equation

$$\mathcal{F}_B = \underbrace{C}_{\text{coordinatized } \Psi^{-1}} \cdot \mathcal{F}_A \cdot \underbrace{E}_{\text{coordinatized } \phi}$$

Challenge: coordinatize mappings ϕ and Ψ^{-1}

Example: DFT

- Consider $\mathcal{A} = \mathbb{C}[x]/x^8 - 1$

$$\mathcal{B} = \langle x^2 \rangle = \mathbb{C}[y]/y^4 - 1, \quad y = x^2$$

- We have

$$\mathbb{C}[y]/y^4 - 1 \quad \xrightarrow[E]{} \quad \mathbb{C}[x]/x^8 - 1$$

$$\text{DFT}_4 \downarrow$$

$$\downarrow \text{DFT}_8$$

$$\bigoplus \mathbb{C}[y]/(y - \omega_4^i) \quad \xleftarrow[C]{\psi^{-1}} \quad \bigoplus \mathbb{C}[x]/(x - \omega_8^i)$$

$$E = \begin{pmatrix} 1 & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot \\ \cdot & 1 & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & 1 & \cdot \\ \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & 1 \\ \cdot & \cdot & \cdot & \cdot \end{pmatrix}$$

$$C = 1/2 \begin{pmatrix} 1 & \cdot & \cdot & \cdot & 1 & \cdot & \cdot & \cdot \\ \cdot & 1 & \cdot & \cdot & \cdot & 1 & \cdot & \cdot \\ \cdot & \cdot & 1 & \cdot & \cdot & \cdot & 1 & \cdot \\ \cdot & \cdot & \cdot & 1 & \cdot & \cdot & \cdot & 1 \end{pmatrix}$$

$$C_2 = \begin{pmatrix} 1 & \cdot \\ \cdot & \cdot & 1 & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & 1 & \cdot & \cdot & \cdot \\ \cdot & 1 \end{pmatrix}$$

$$\text{DFT}_4 = C \cdot \text{DFT}_8 \cdot E$$

$$E_2 = \begin{pmatrix} 1 & \cdot & \cdot & \cdot \\ \cdot & 1 & \cdot & \cdot \\ \cdot & \cdot & 1 & \cdot \\ \cdot & \cdot & \cdot & 1 \end{pmatrix}$$

Crucial parameters: basis in $\mathcal{A}, \mathcal{B}, \mathcal{A}_i, \mathcal{B}_i$

Multiple degrees of freedom

Example: DCT-3

- Consider $\mathcal{A} = \mathbb{R}[x]/x^8 + 1$

$$\mathcal{B} = \mathbb{R}[y]/T_4(y), \quad y = (x + x^{-1})/2$$

- We have

$$\mathbb{R}[y]/T_4(y) \xrightarrow[E]{} \mathbb{R}[x]/x^8 + 1$$

$$\text{DCT-3}_4 \downarrow$$

$$\downarrow \text{RDFT-3}_8$$

$$\bigoplus \mathbb{C}[y]/(y - \cos \alpha_i) \xleftarrow[C]{\psi^{-1}} \bigoplus \mathbb{C}[x]/(x - \omega_8^i)(x - \omega_8^{8-i})$$

$$E = W_A = \begin{pmatrix} 1 & \cdot & \cdot & \cdot \\ \cdot & 1/2 & \cdot & \cdot \\ \cdot & \cdot & 1/2 & \cdot \\ \cdot & \cdot & \cdot & 1/2 \\ \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & -1/2 \\ \cdot & \cdot & -1/2 & \cdot \\ \cdot & -1/2 & \cdot & \cdot \end{pmatrix}$$

$$C = DS_1 = \begin{pmatrix} 1 & \cdot \\ \cdot & \cdot & 1 & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & 1 & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & 1 & \cdot \end{pmatrix}$$

$$\text{DCT-3}_4 = DS_1 \cdot \text{RDFT-3}_8 \cdot W_A$$

$$\text{DST-3}_4 = DS_2 \cdot \text{RDFT-3}_8 \cdot W_{S0}$$

Outline

1. Introduction
2. Embedding method
3. **Algorithm Embedding method.**

Projection Method: Problem statement

■ Given:

- Regular signal model $(\mathcal{A}, \Phi_{\mathcal{A}})$ + Fourier transform $\mathcal{F}_{\mathcal{A}}$
- Cooley-Tukey type algorithm for $\mathcal{F}_{\mathcal{A}}$ made with sub-algebra \mathcal{C}

- Regular signal model $(\mathcal{B}, \Phi_{\mathcal{B}})$ + Fourier transform $\mathcal{F}_{\mathcal{B}}$
- $\mathcal{B} \leq \mathcal{A}$

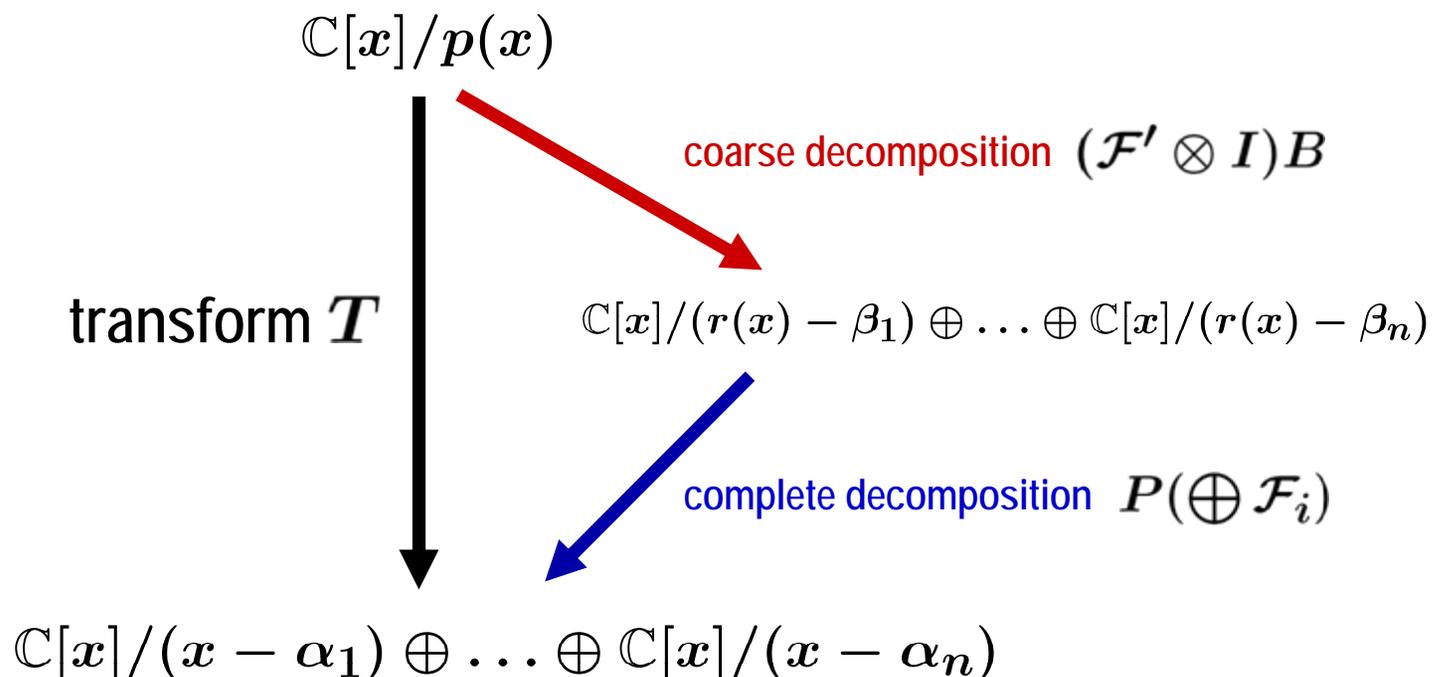
■ Find:

- Cooley-Tukey type algorithm for $\mathcal{F}_{\mathcal{B}}$

Reminder: Cooley-Tukey Type Algorithms

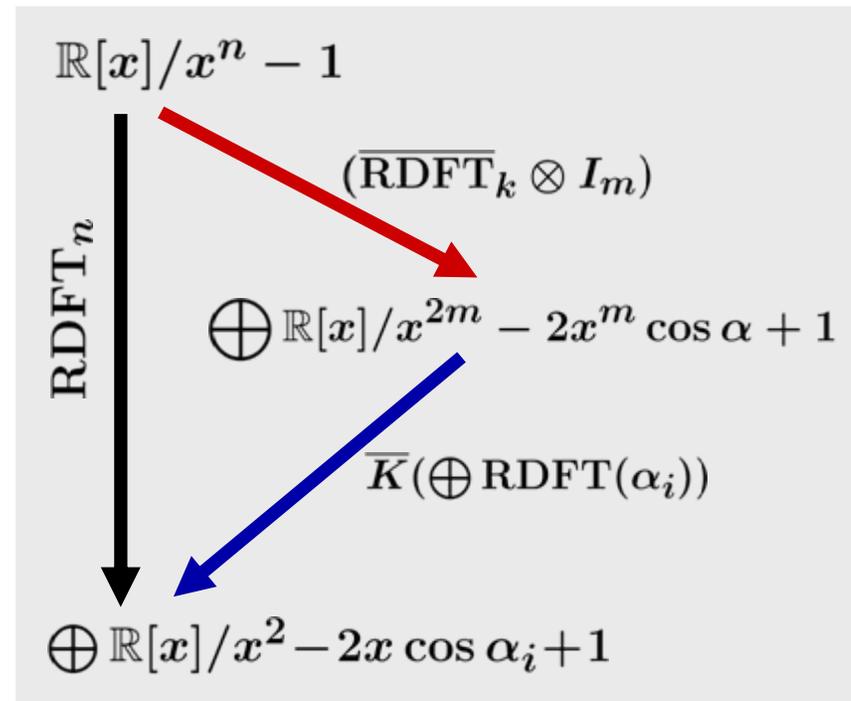
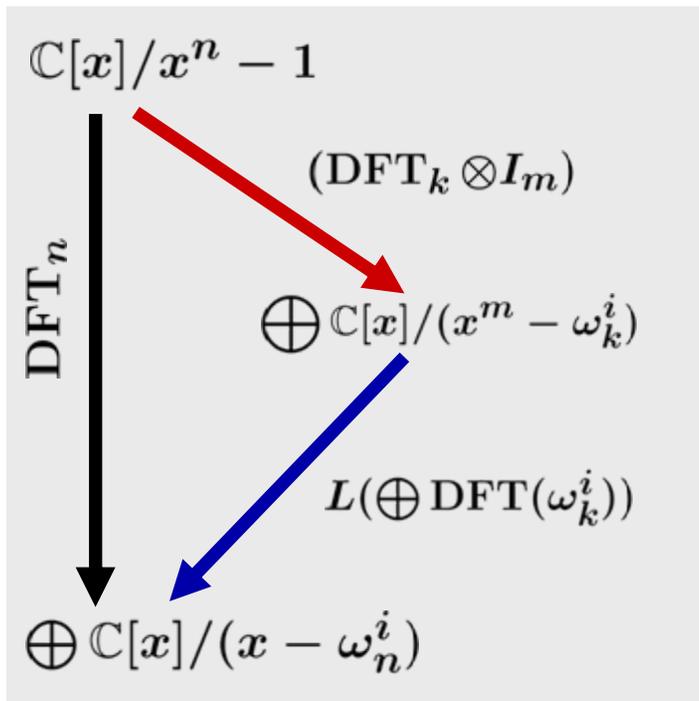
Assume

$$p(x) = q(r(x))$$



Fast Algorithms DFT and RDFT

Decomposition: $x^n - 1 = (x^m)^k - 1$

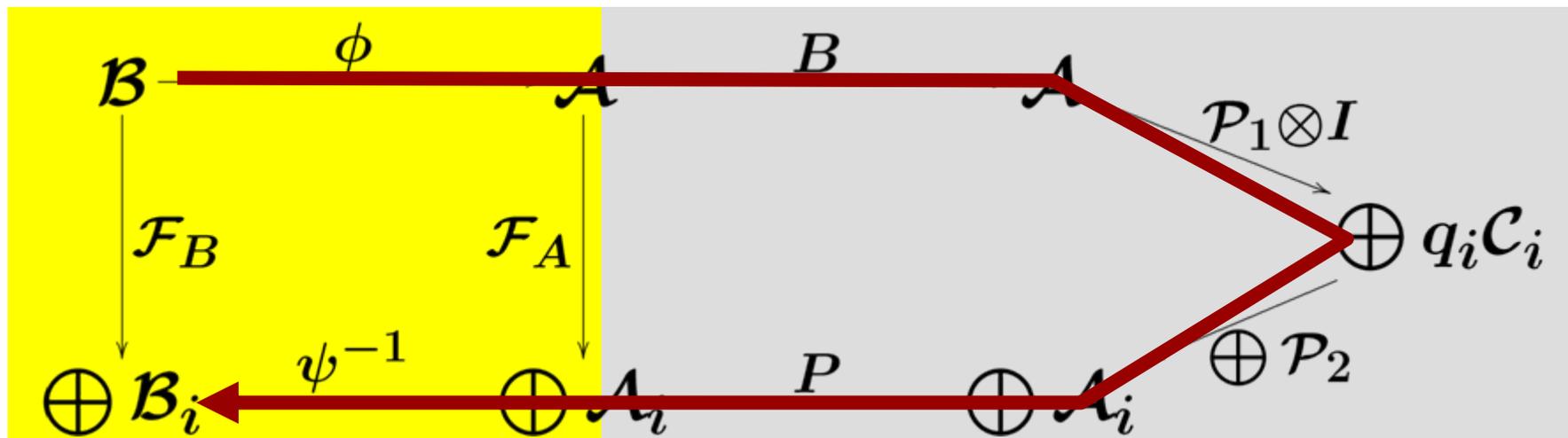


$$\text{DFT}_n = \underline{L_m^n (I_k \otimes \text{DFT}_m)} \underline{T_m^n (\text{DFT}_k \otimes I_m)}$$

$$\text{RDFT}_n = \underline{(K_{m/2}^{n/2} \otimes I_2) (\oplus \text{RDFT}_m(\alpha_i))} \underline{(\overline{\text{RDFT}}_k \otimes I_m)}$$

Algorithm Projection Method

- Basic idea: embedding and Cooley-Tukey
- Following diagram commutes



The embedding

The \mathcal{F}_A Cooley-Tukey

Matrix equation

$$\mathcal{F}_B = C \cdot P \cdot (\bigoplus P_2) \cdot (P_1 \otimes I) \cdot B \cdot E$$

coordinatized ψ^{-1} spectral permutation fine decomposition coarse decomposition base change coordinatized ϕ

Challenge: choose ϕ and ψ^{-1} for sparse factorization

Algorithm Projection Steps

- Algorithm derivation is mechanical rewriting of

$$\mathcal{F}_B \rightarrow \underbrace{C \cdot P\left(\bigoplus \mathcal{P}_2(\beta_i)\right)}_{\text{apply left rules}} \underbrace{(\mathcal{P}_1 \otimes I) B \cdot E}_{\text{apply right rules}}$$

$$\mathcal{F}_B = C \cdot \mathcal{F}_A \cdot E$$

$$C \cdot \mathcal{F}_A \rightarrow E^{-1} \cdot \mathcal{F}_B \quad \mathcal{F}_A \cdot E \rightarrow C^{-1} \cdot \mathcal{F}_B$$

left "push" rules
right "pull" rules

- Step 1: Write down the embedding
- Step 2: Coordinatize and obtain right / left rules
- Step 3: Plug in \mathcal{F}_A Cooley-Tukey and apply rules

But remember that E and C are chosen!

Example: Embedding DFT_4 into DFT_8

- Step 1: Write down the embedding

$$\mathcal{A} = \mathbb{C}[x]/x^8 - 1$$

$$\mathcal{B} = \mathbb{C}[y]/y^4 - 1 = \langle x^2 \rangle$$

- Step 2: Coordinatize, and obtain right/left rules

$$C_n \cdot DFT_n \cdot E_n = DFT_{n/2},$$

$$DFT_n E_n = C_n^{-1} DFT_{n/2},$$

$$C_n DFT_n = DFT_{n/2} E_n^{-1}.$$

$$C_8 = 1/2 \begin{pmatrix} 1 & \cdot & \cdot & \cdot & 1 & \cdot & \cdot & \cdot \\ \cdot & 1 & \cdot & \cdot & \cdot & 1 & \cdot & \cdot \\ \cdot & \cdot & 1 & \cdot & \cdot & \cdot & 1 & \cdot \\ \cdot & \cdot & \cdot & 1 & \cdot & \cdot & \cdot & 1 \end{pmatrix}$$

$$E_8 = \begin{pmatrix} 1 & \cdot & \cdot & \cdot \\ \cdot & 1 & \cdot & \cdot \\ \cdot & \cdot & 1 & \cdot \\ \cdot & \cdot & \cdot & 1 \\ \cdot & \cdot & \cdot & \cdot \end{pmatrix}$$

- Step 3: Plug in Cooley-Tukey for DFT_8 , apply rules

$$DFT_4 = C_8 \cdot L_4^8(\oplus DFT_4(r_i)) (DFT_2 \otimes I_4) \cdot E_8$$

$$DFT_4 = C_8 \cdot L_4^8(\oplus DFT_4(r_i)) E_8 (DFT_2 \otimes I_2)$$

$$DFT_4 = L_4^8(\oplus C_4 DFT_4(r_i)) E_8 (DFT_2 \otimes I_2)$$

$$DFT_4 = L_4^8(\oplus DFT_2(r_i)) (DFT_2 \otimes I_2)$$

Example: Embed DCT-3₄ into RDFT-3₈

- Step 1: Write down the embedding

$$\mathcal{A} = \mathbb{R}[x]/x^8 + 1$$

$$\mathcal{B} = \mathbb{R}[y]/T_4(y) = \langle (x + x^{-1})/2 \rangle$$

- Step 2: Coordinatize, and obtain right/left rules

$$C_n \cdot \text{RDFT-3}_n \cdot E_n = \text{DCT-3}_{n/2},$$

$$\text{RDFT-3}_n E_n = C_n^{-1} \text{DCT-3}_{n/2},$$

$$C_n \text{RDFT-3}_n = \text{DCT-3}_{n/2} E_n^{-1}.$$

$$C_8 = \begin{pmatrix} 1 & \cdot \\ \cdot & \cdot & 1 & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & 1 & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & 1 & \cdot \end{pmatrix}$$

$$E_8 = \frac{1}{2} \begin{pmatrix} 2 & \cdot & \cdot & \cdot \\ \cdot & 1 & \cdot & \cdot \\ \cdot & \cdot & 1 & \cdot \\ \cdot & \cdot & \cdot & 1 \\ \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & -1 \\ \cdot & \cdot & -1 & \cdot \\ \cdot & -1 & \cdot & \cdot \end{pmatrix}$$

- Step 3: Plug in Cooley-Tukey for RDFT-3₈, apply rules

$$\text{DCT-3}_4 = C_8 \cdot (K_2^4 \otimes I_2) \left(\bigoplus \text{RDFT-3}_4(r_i) \right) (\overline{\text{RDFT-3}_4} \otimes I_2) \cdot E_8$$

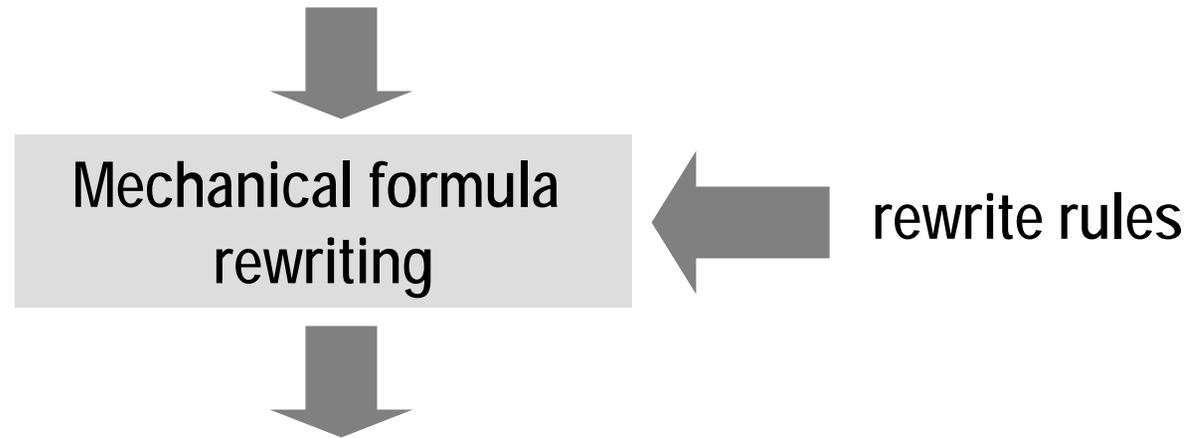
$$\text{DCT-3}_4 = \underbrace{K_2^4}_{\text{permutations}} \left(\bigoplus \text{DCT-3}_2(r_i) \right) \underbrace{Q}_{\text{permutations}} (\overline{\text{DCT-3}} \oplus (I \otimes \text{RDFT-3})) \underbrace{G}_{\text{permutations}} \overline{\text{DCT-4}} \underbrace{P}_{\text{permutations}}$$

permutations

Conclusion :

Automatic Derivation of Algorithms

$$\text{RDFT}_n = (K_{m/2}^{n/2} \otimes I_2) (\bigoplus \text{RDFT}_m(\alpha_i)) (\overline{\text{RDFT}_k} \otimes I_m)$$



$$\text{DCT-3}_k = K_m^{mn} (\bigoplus \text{DCT-3}_m(2r_i)) Q (\overline{\text{DCT-3}_n} \oplus (I_{m/2-1} \otimes \text{RDFT-3}_{2n}) G \oplus \overline{\text{DCT-4}_n}) R (L_m^{2nm})$$

- Number of rewrite rules is small
- Hard to derive by hand
- Gives new, composite algorithms
- Nice way to handle the ever growing number of transforms

Questions?